

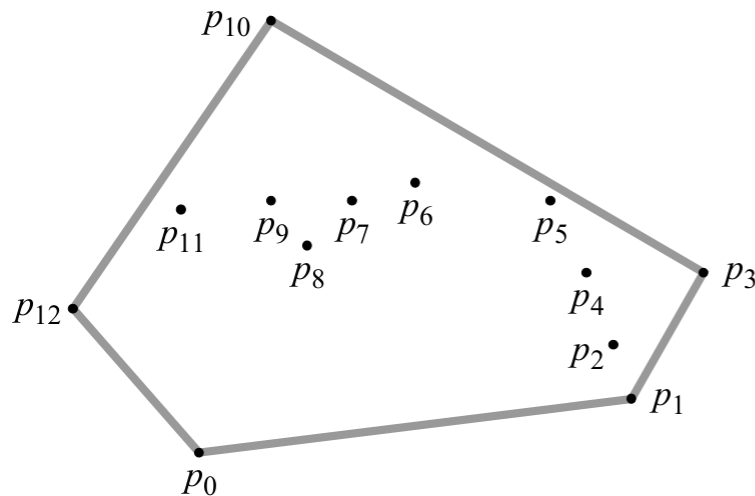
Otoczka wypukła

(ang. Convex Hull)

2026

The convex hull problem

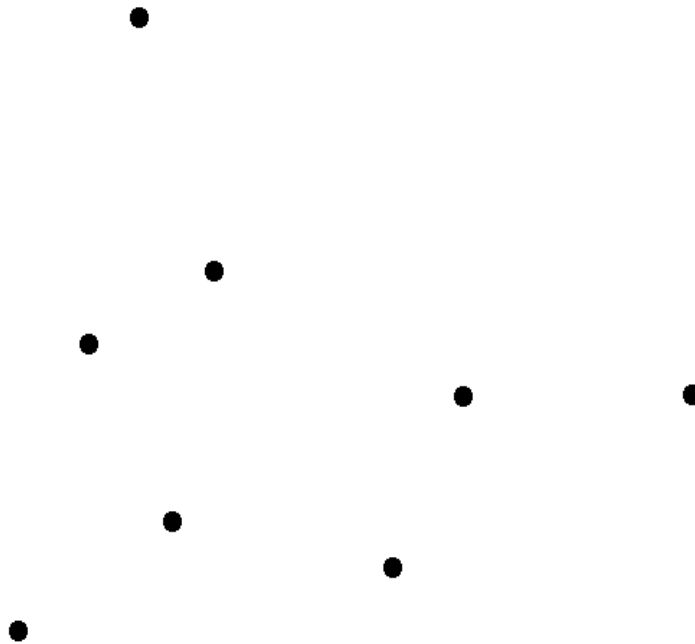
- Otoczka wypukła zbioru punktów Q to najmniejszy wielokąt wypukły P taki, że każdy punkt ze zbioru Q leży albo na brzegu P , albo w jego wnętrzu.



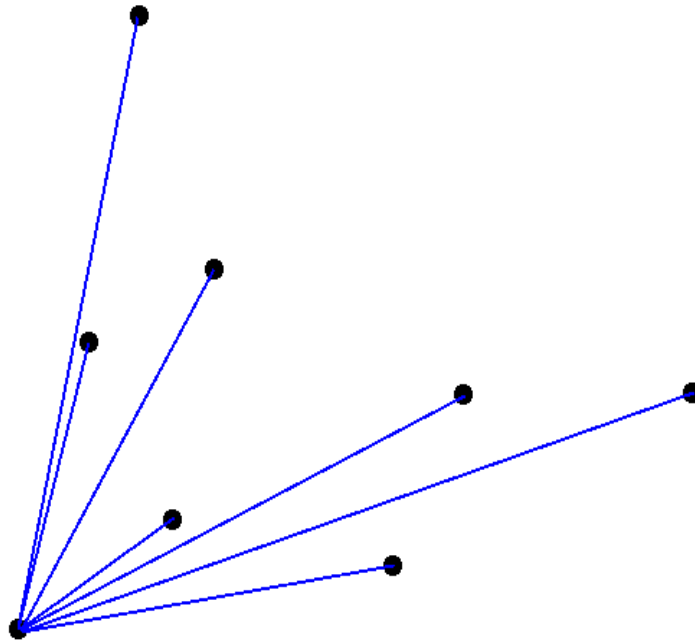
Graham's Scan

- Start at point guaranteed to be on the hull. (the point with the minimum y value)
- **Sort** remaining points by **polar angles** of vertices relative to the first point.
- Go through sorted points, keeping vertices of points that have **left turns** and dropping points that have **right turns**.

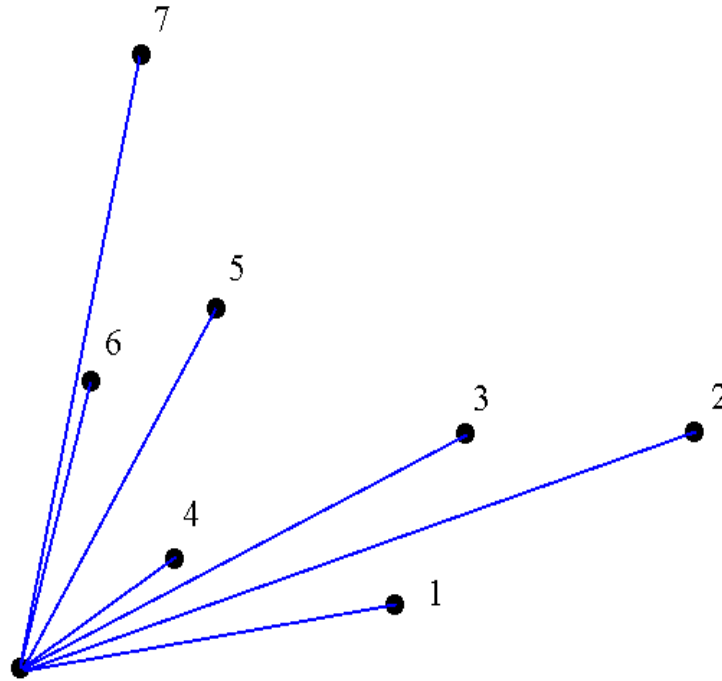
Graham's Scan



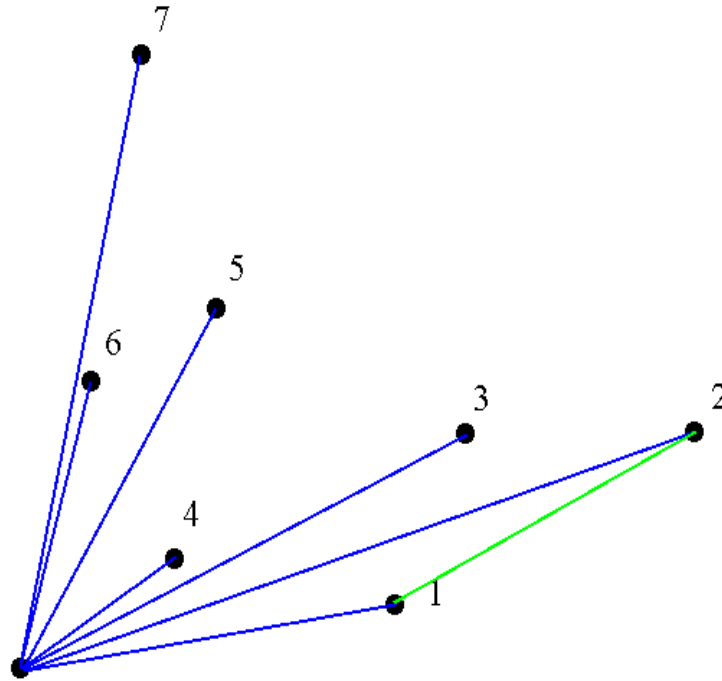
Graham's Scan



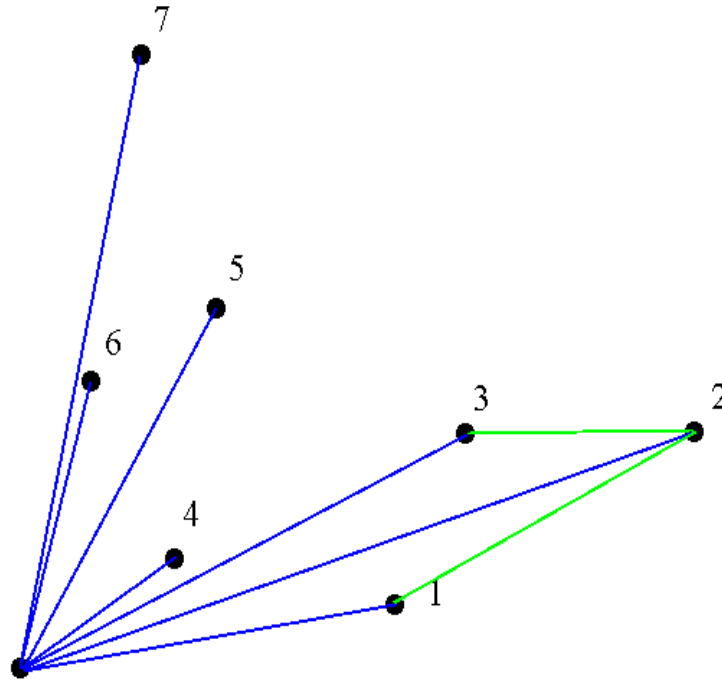
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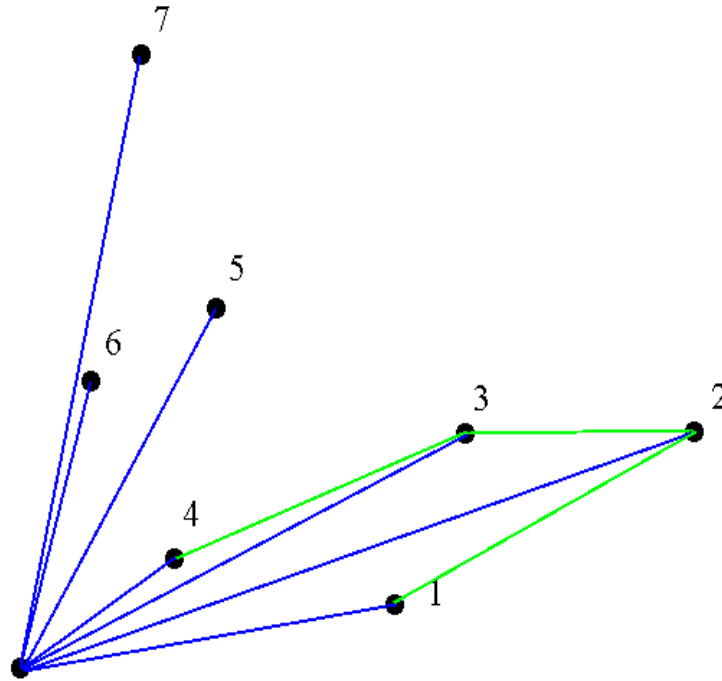
Graham's Scan



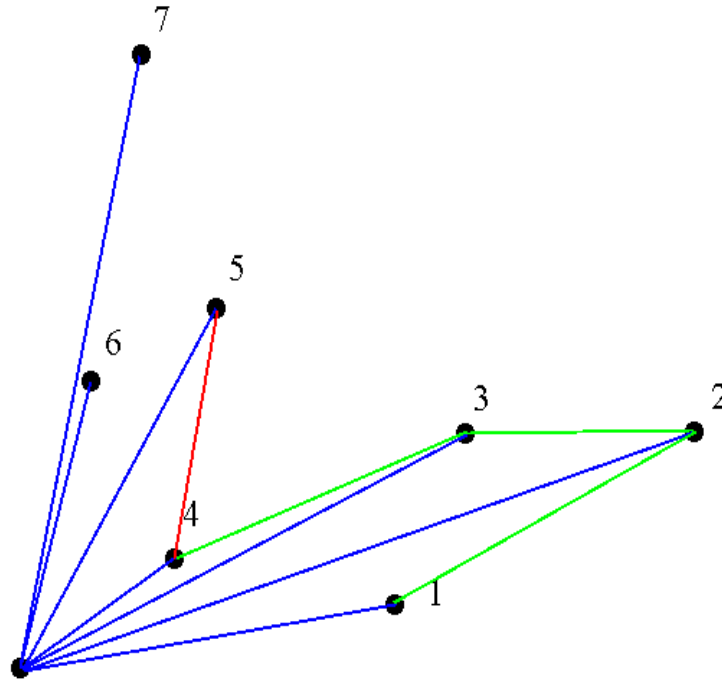
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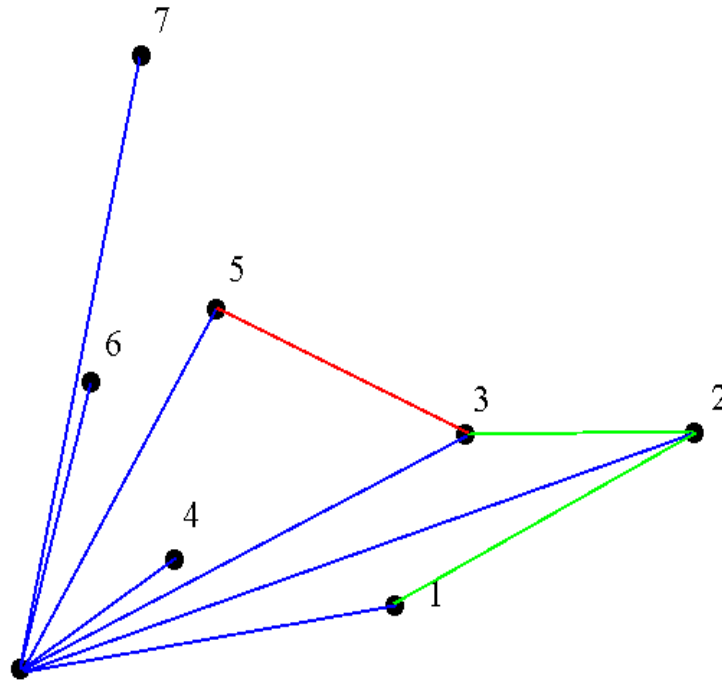
Graham's Scan



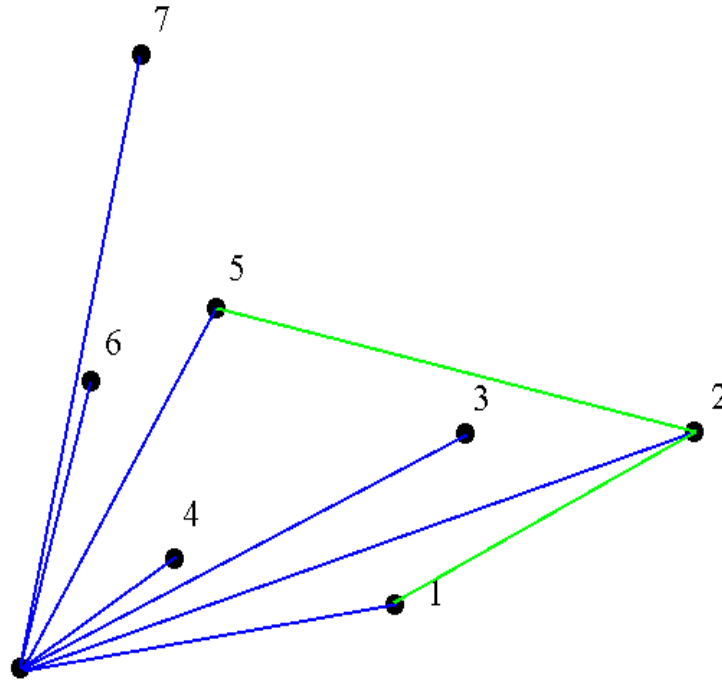
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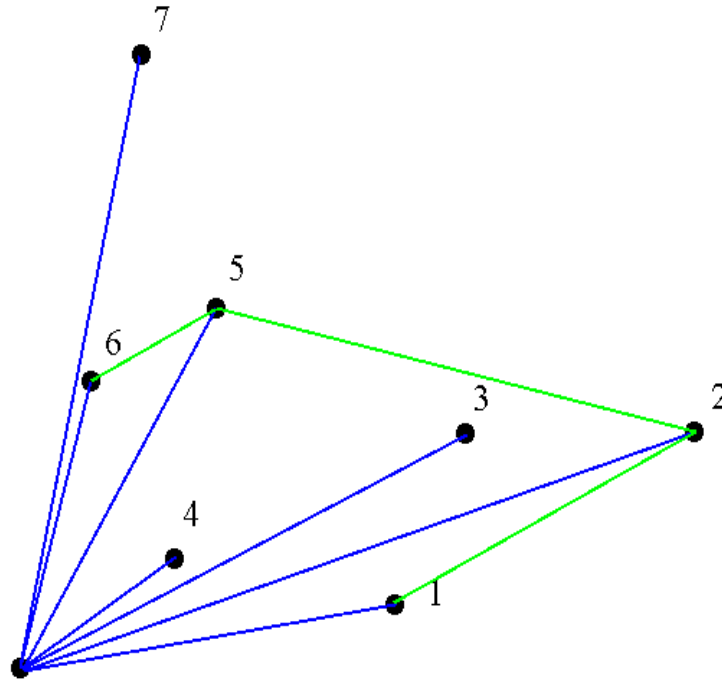
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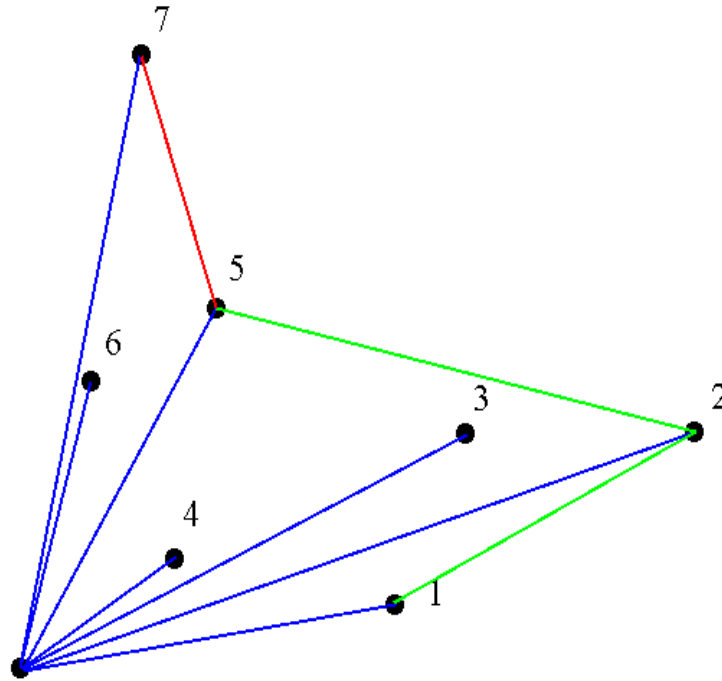
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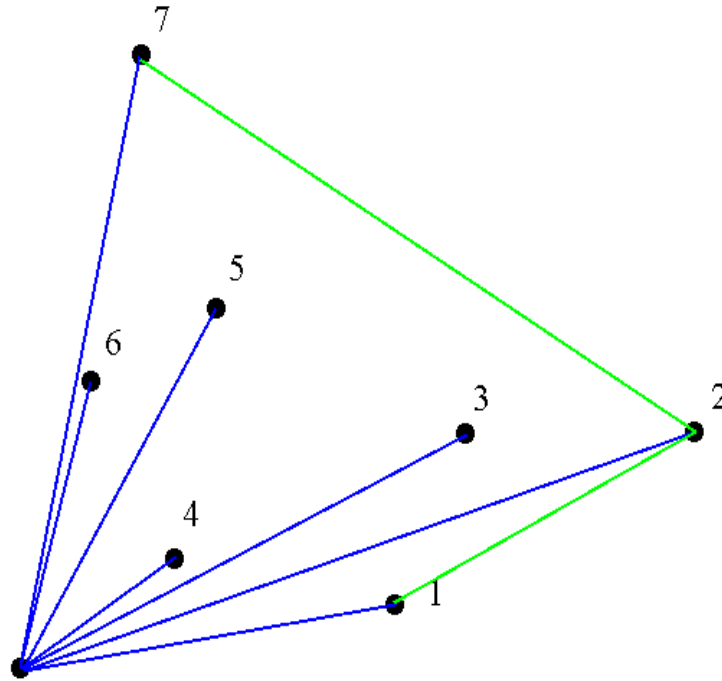
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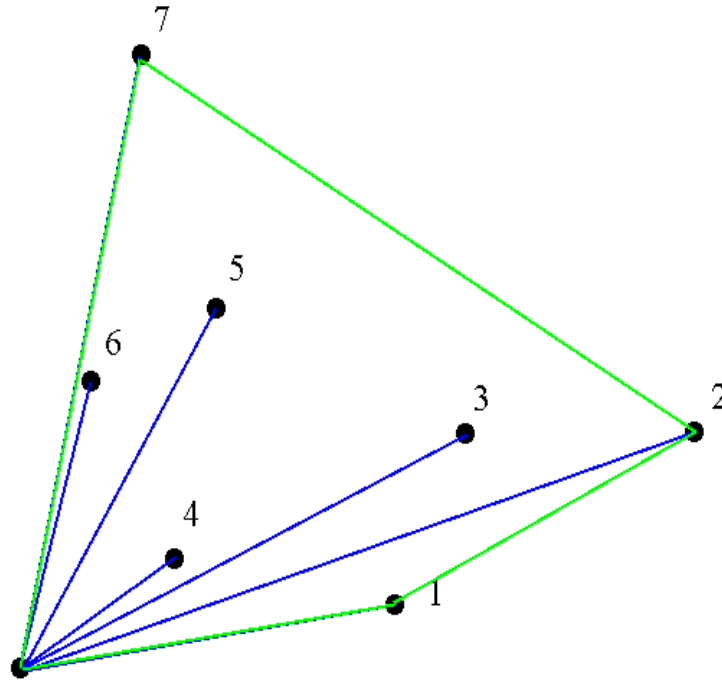
Graham's Scan



Graham's Scan



Graham's Scan



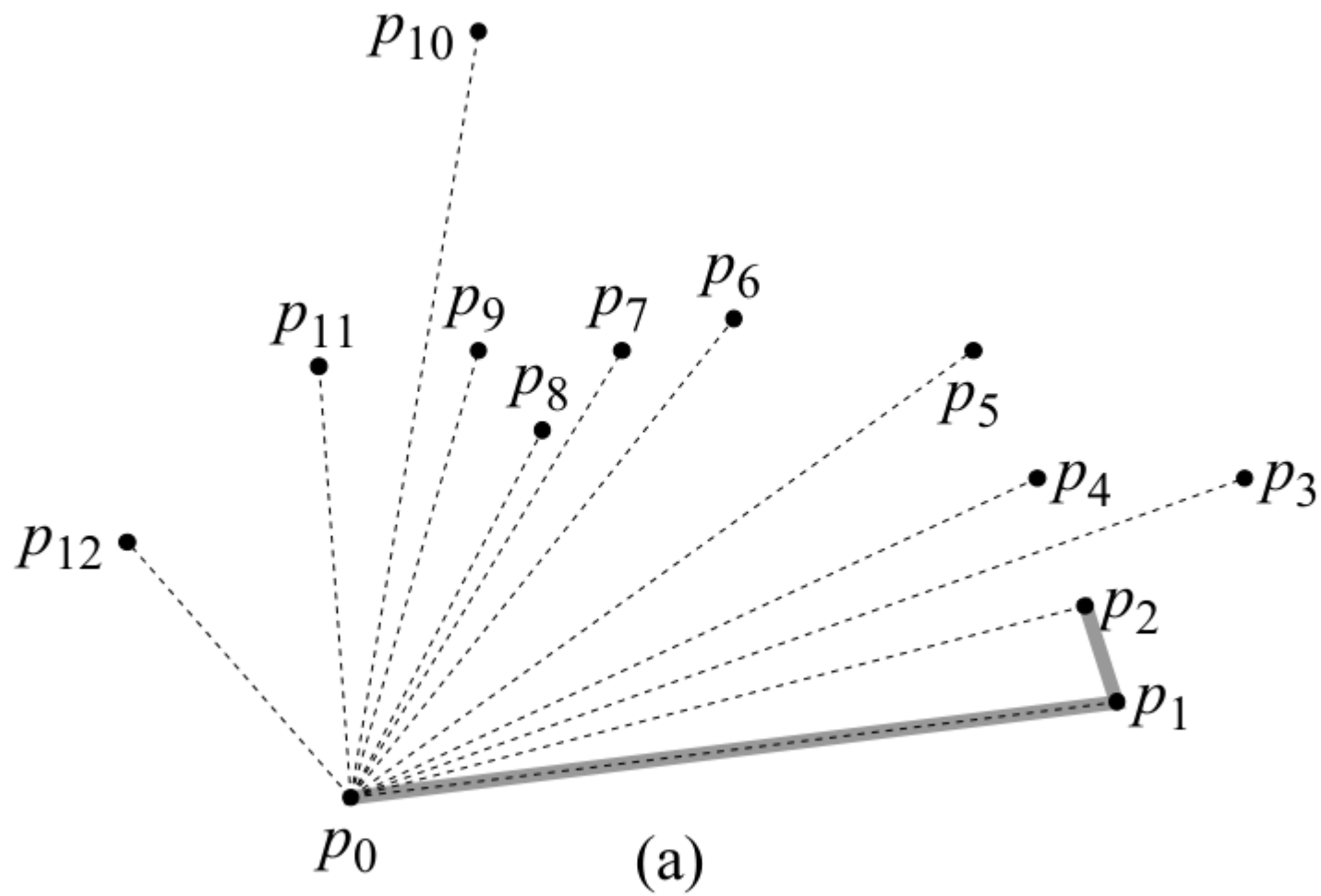
Graham's Runtime

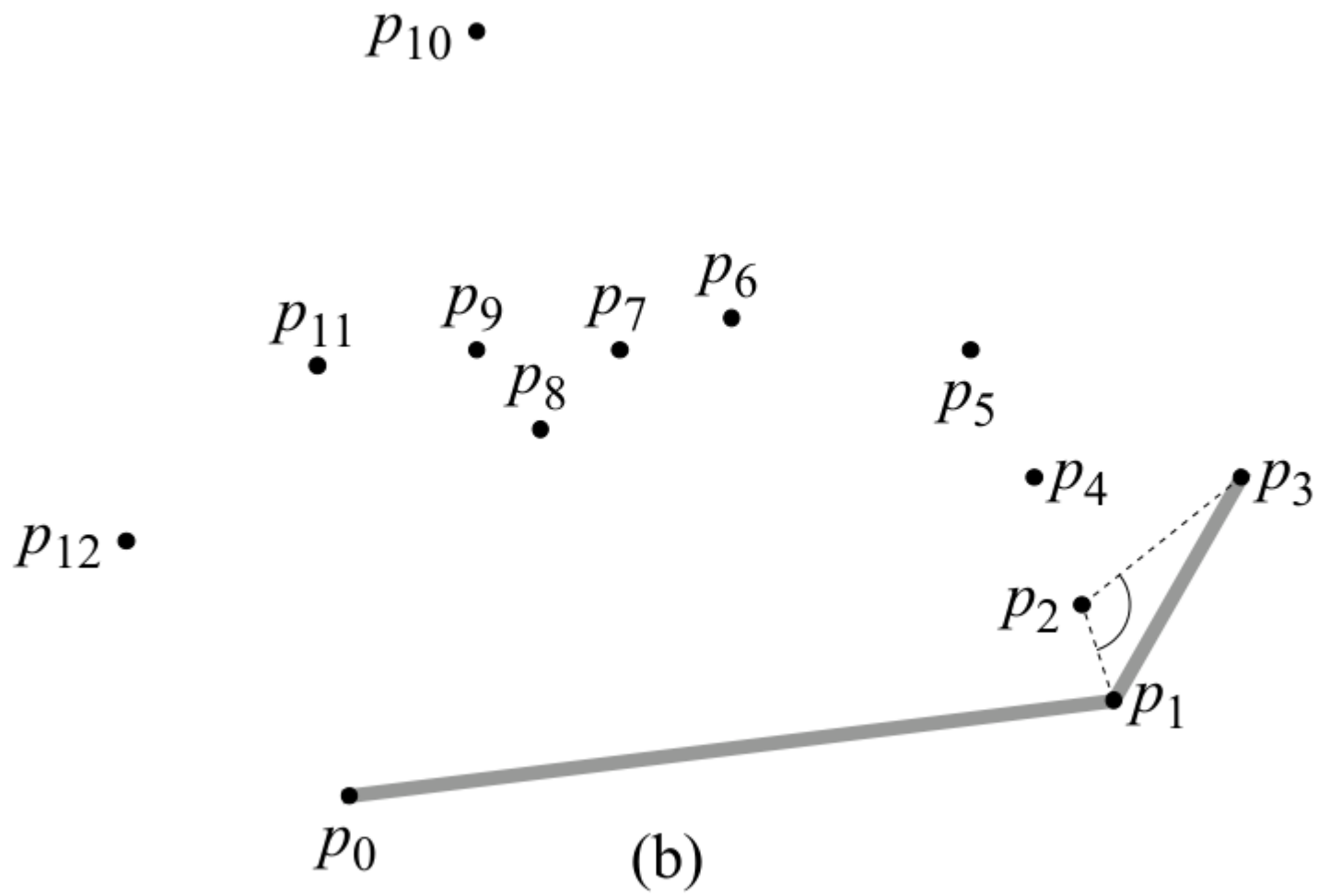
- Graham's scan is $O(n \log n)$ due to initial sort of angles.

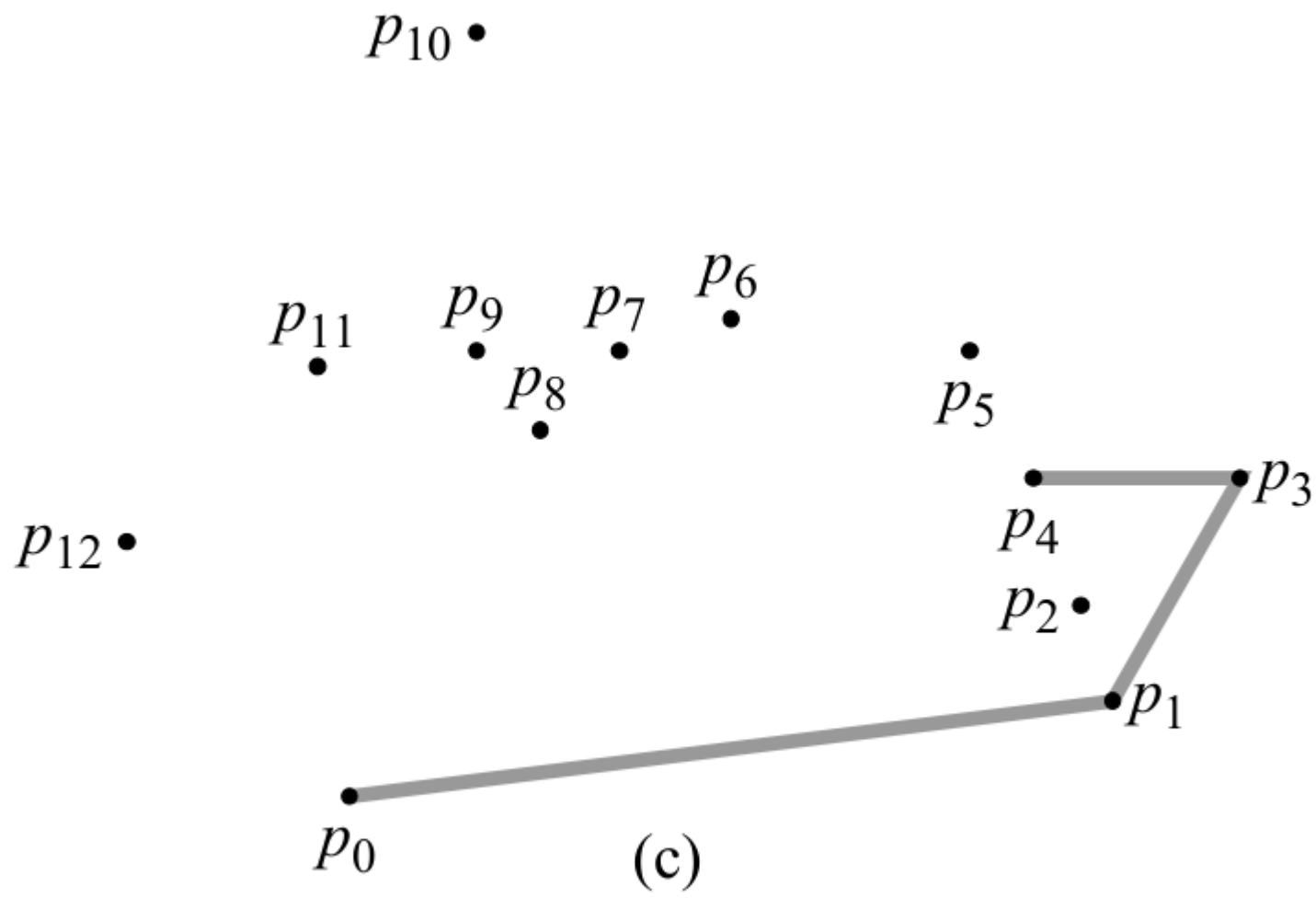
■ A more detailed algorithm

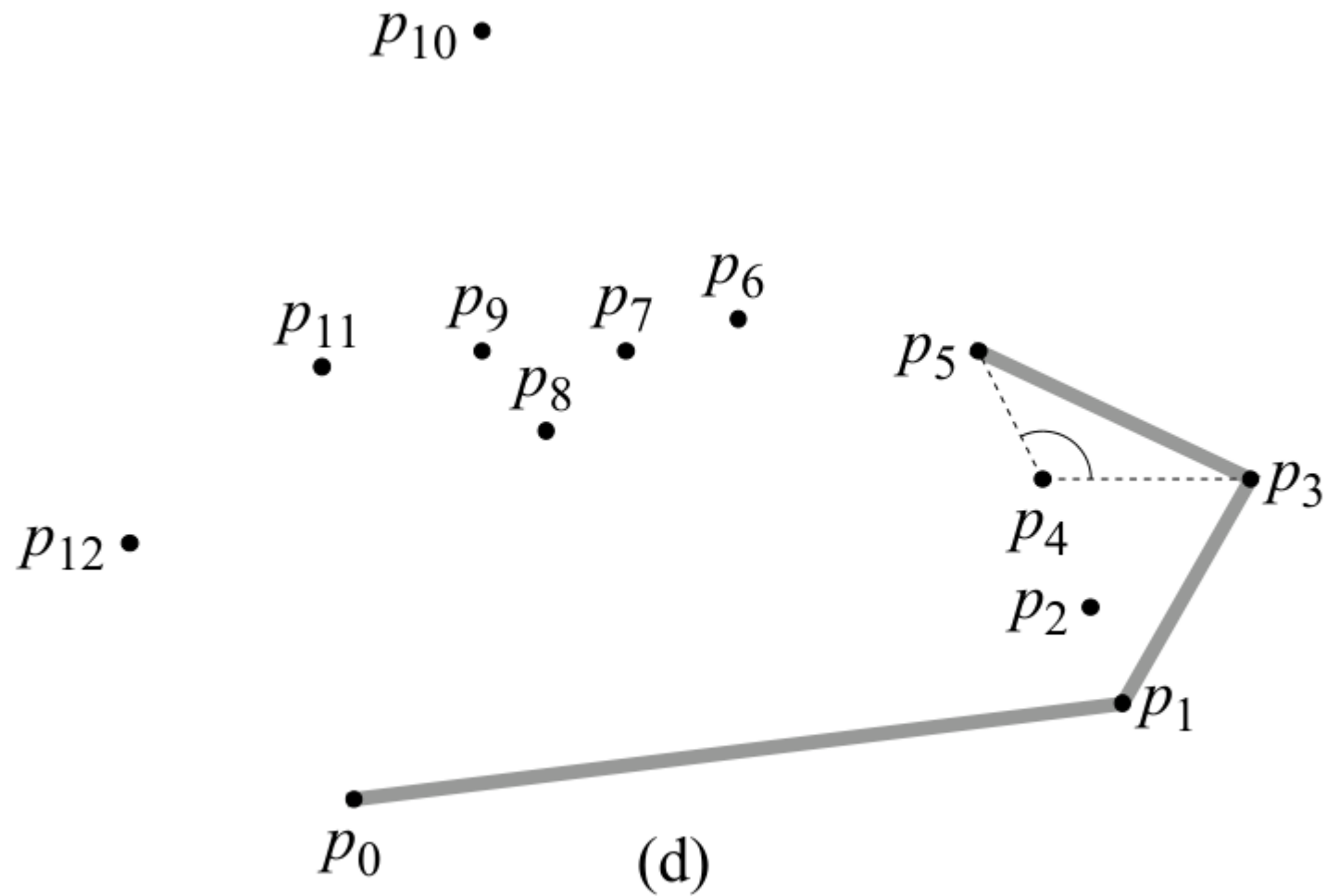
GRAHAM-SCAN(Q)

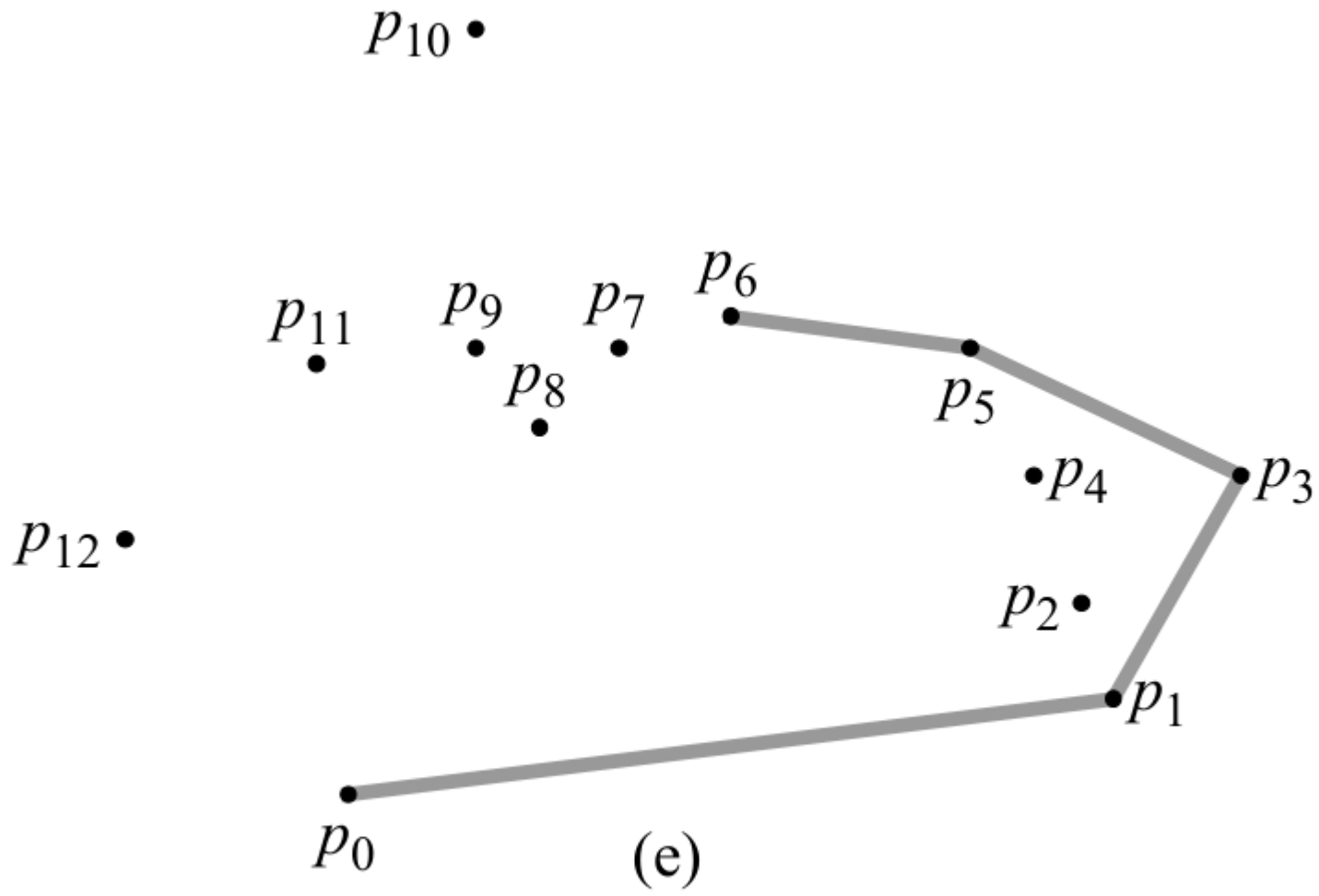
- 1 let p_0 be the point in Q with the minimum y -coordinate,
or the leftmost such point in case of a tie
- 2 let $\langle p_1, p_2, \dots, p_m \rangle$ be the remaining points in Q ,
sorted by polar angle in counterclockwise order around p_0
(if more than one point has the same angle, remove all but
the one that is farthest from p_0)
- 3 PUSH(p_0, S)
- 4 PUSH(p_1, S)
- 5 PUSH(p_2, S)
- 6 **for** $i \leftarrow 3$ **to** m
- 7 **do while** the angle formed by points NEXT-TO-TOP(S), TOP(S),
 and p_i makes a nonleft turn
- 8 **do** POP(S)
- 9 PUSH(p_i, S)
- 10 **return** S



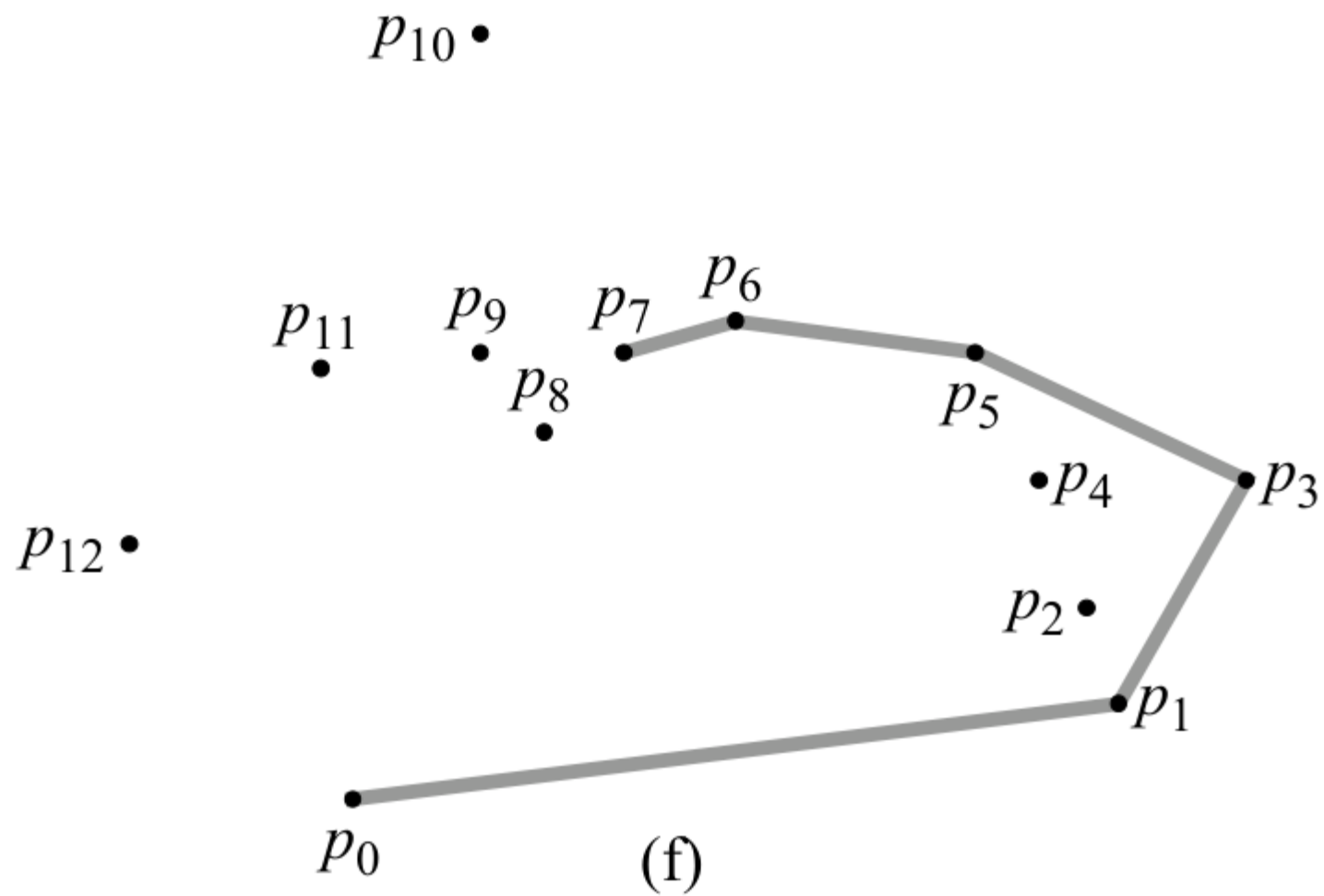


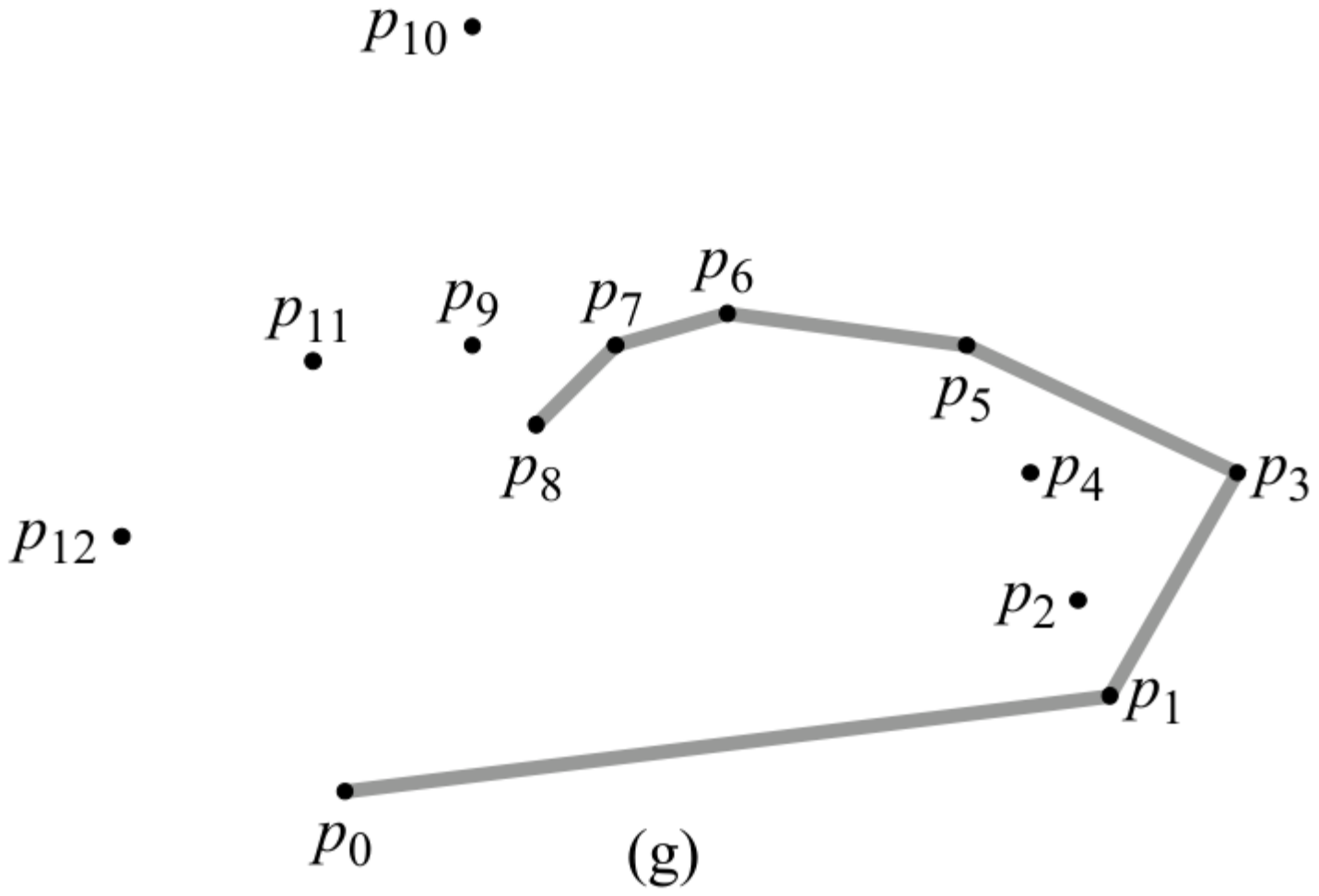


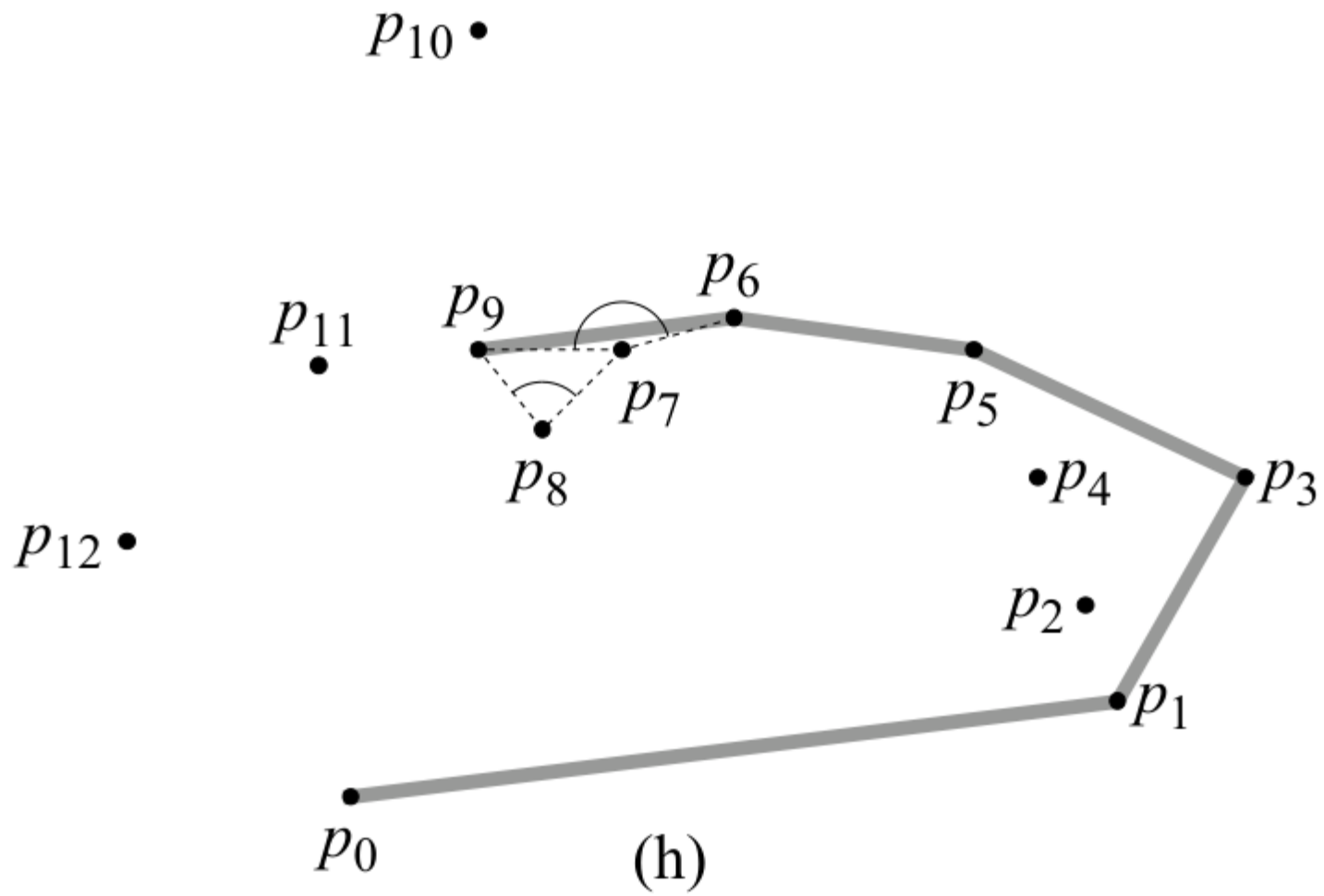


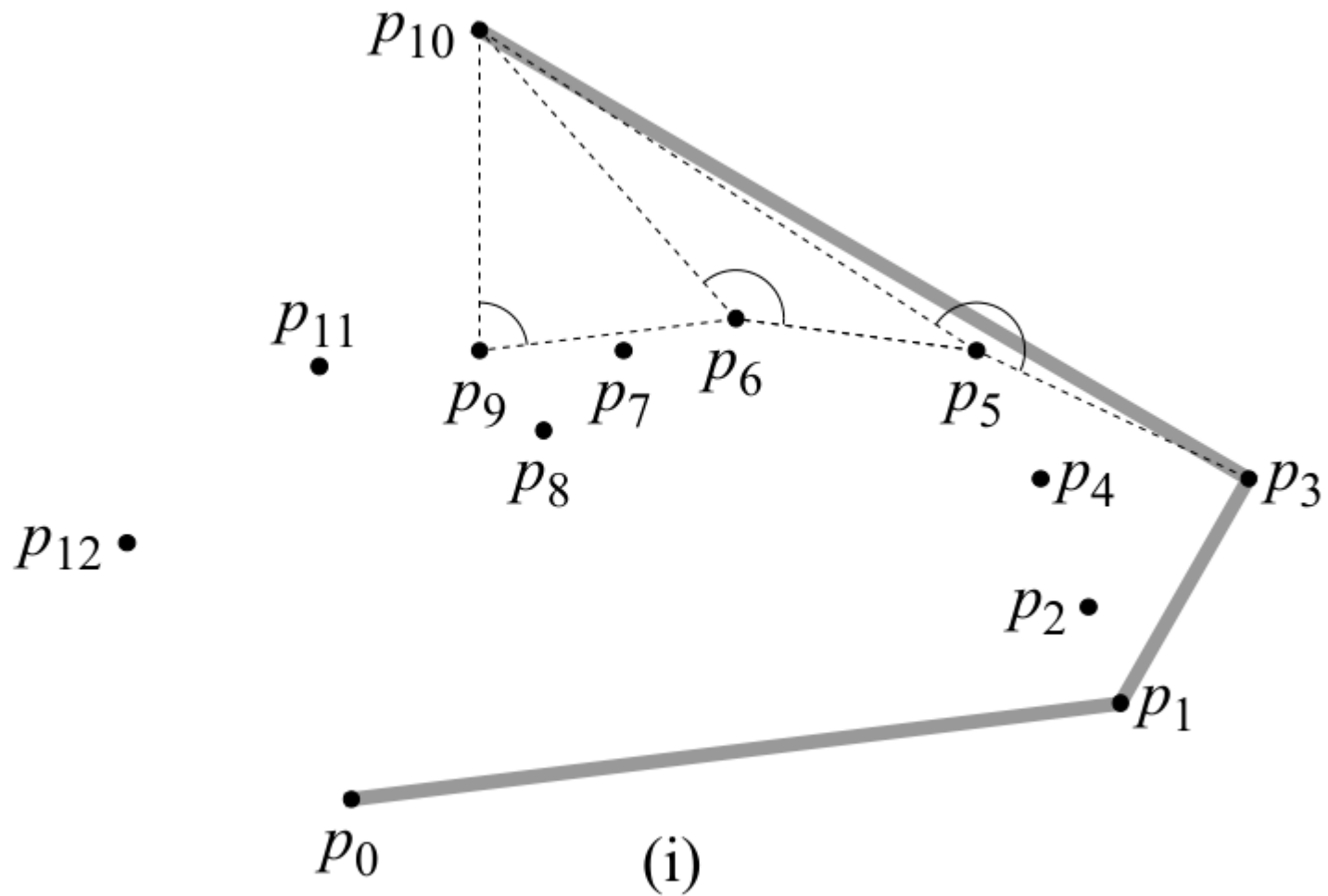


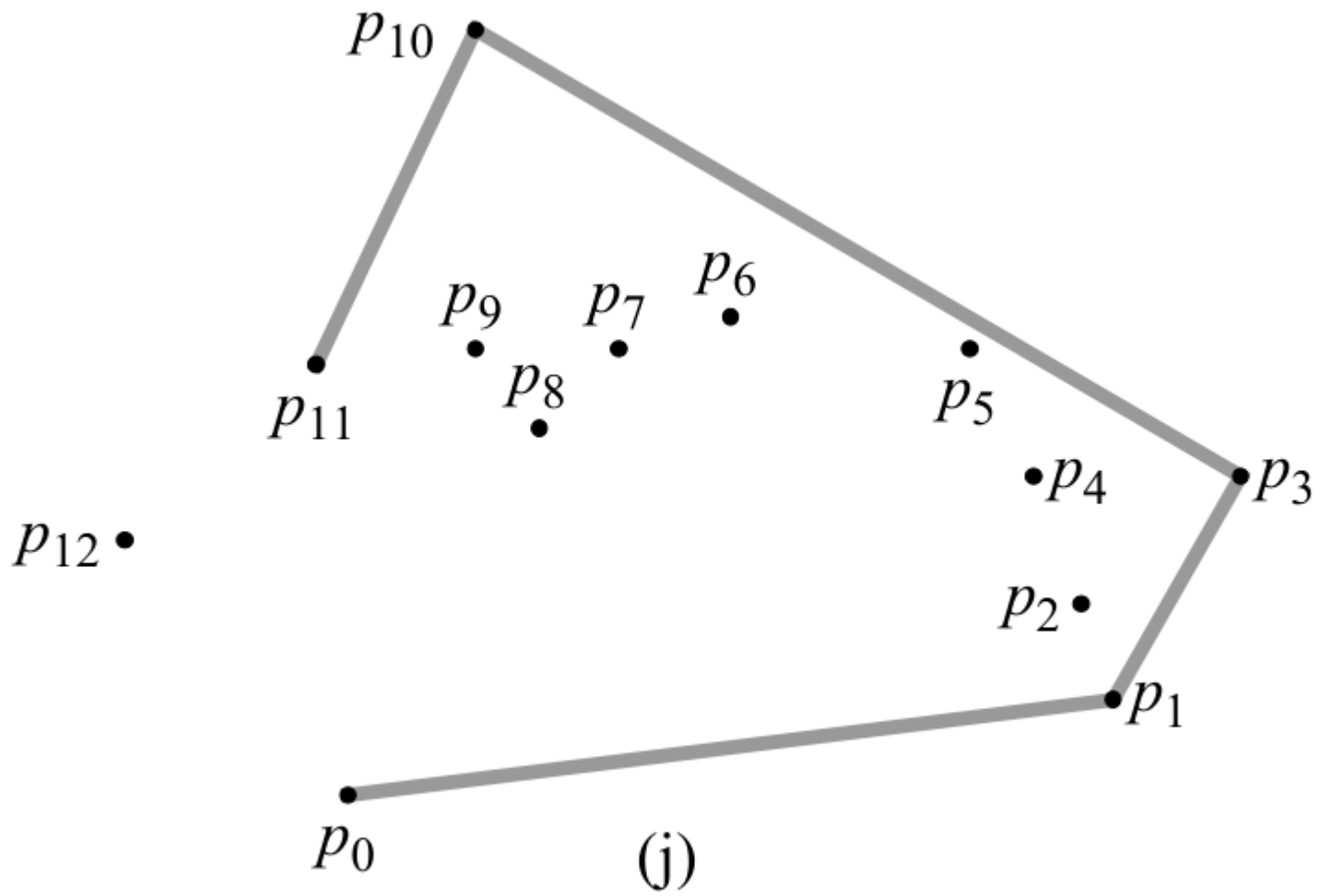
(e)

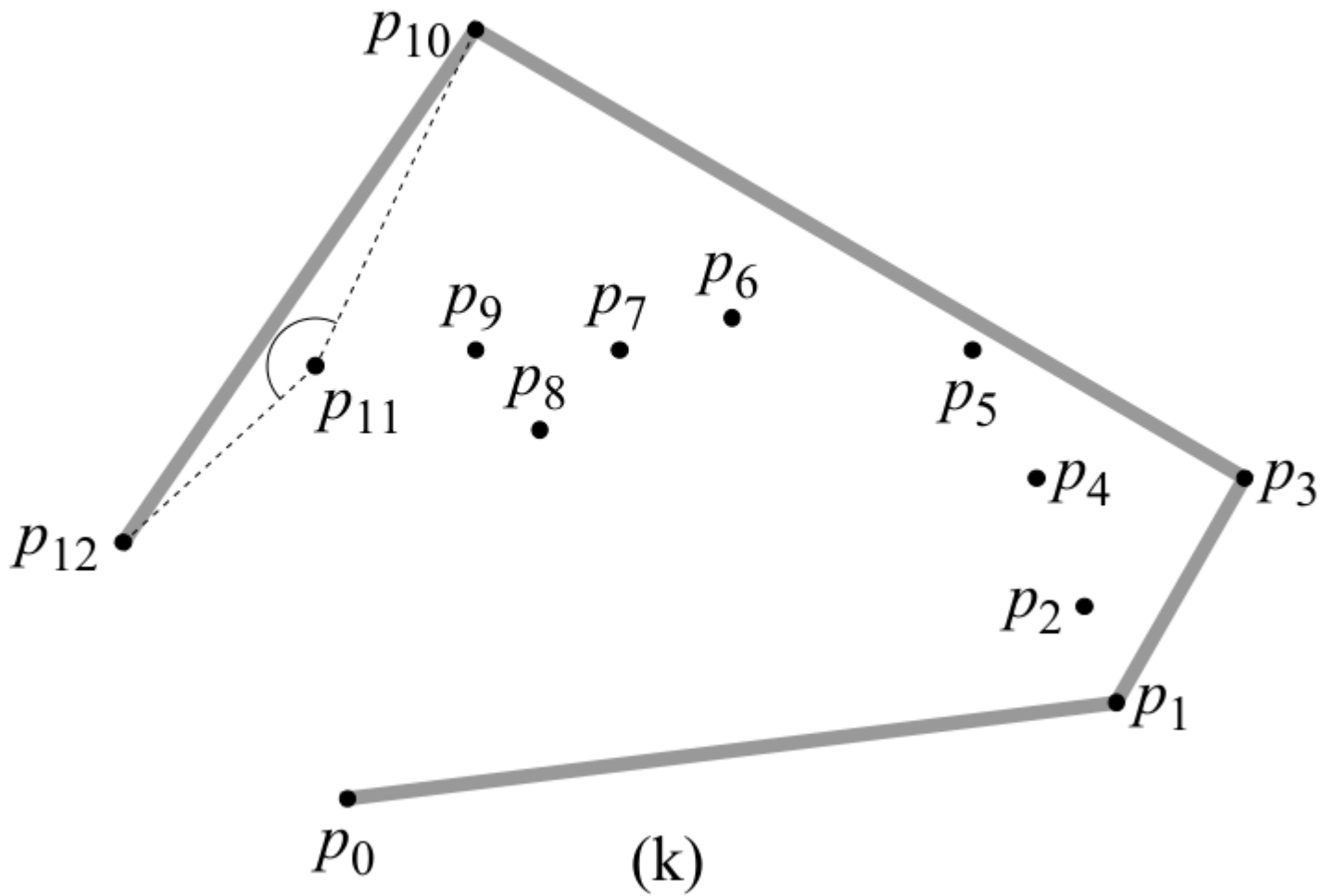


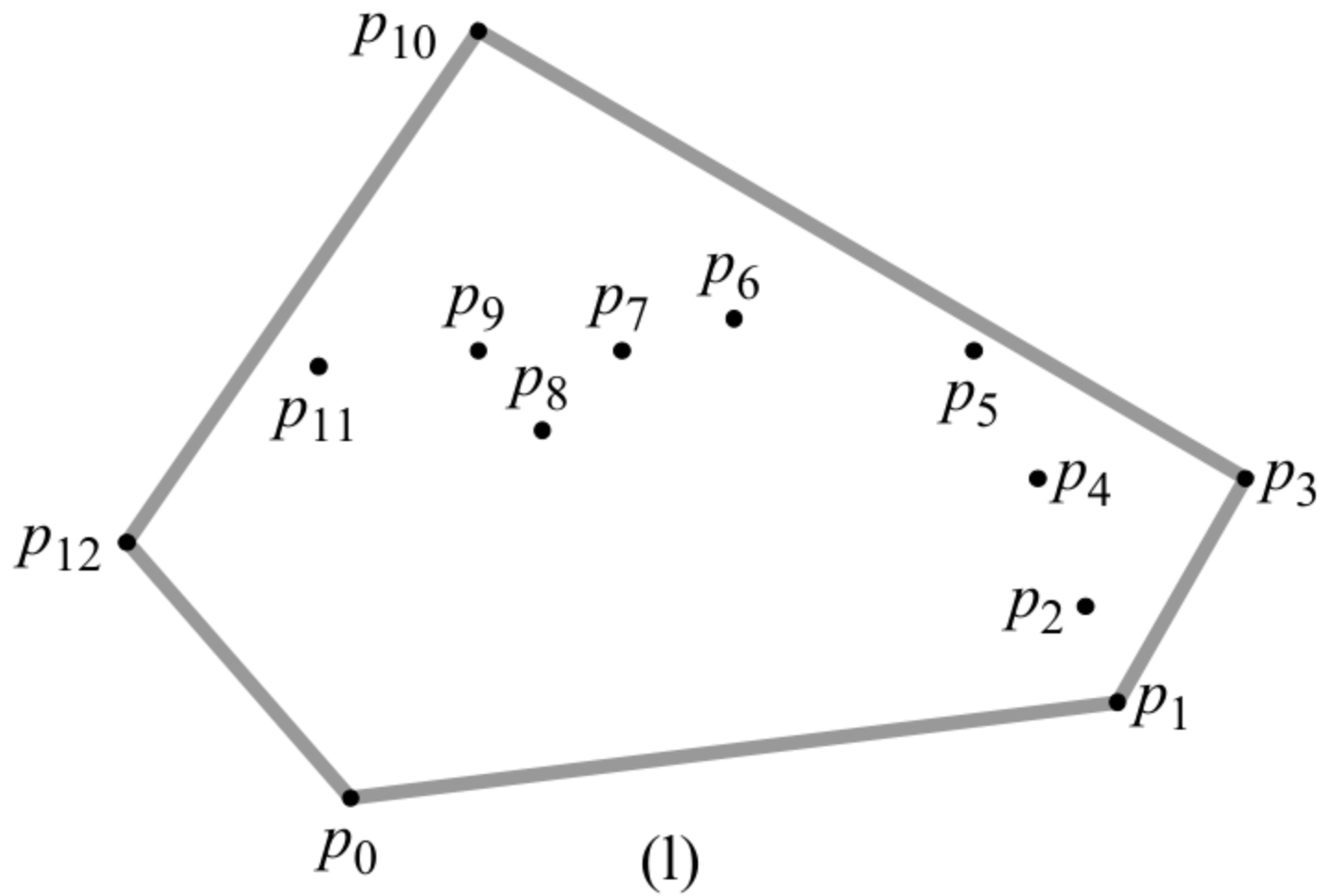






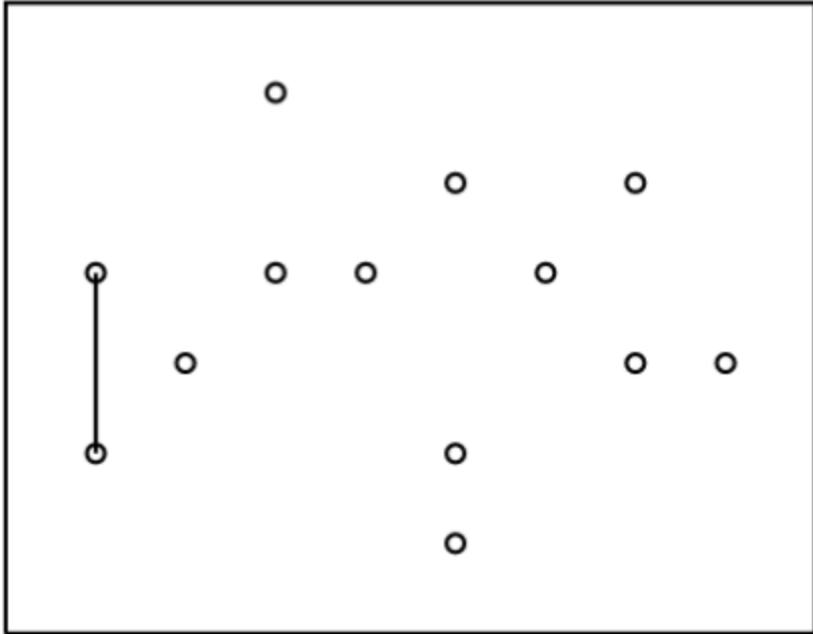


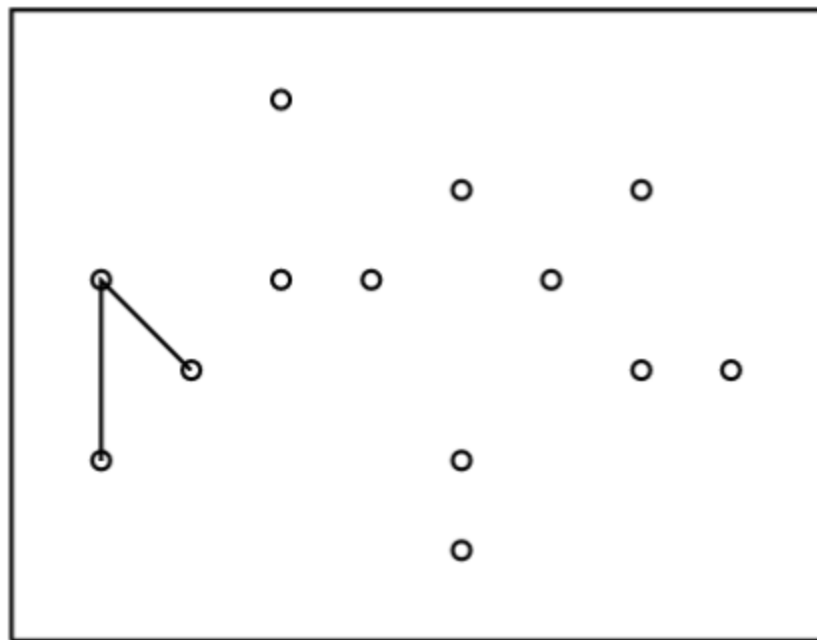


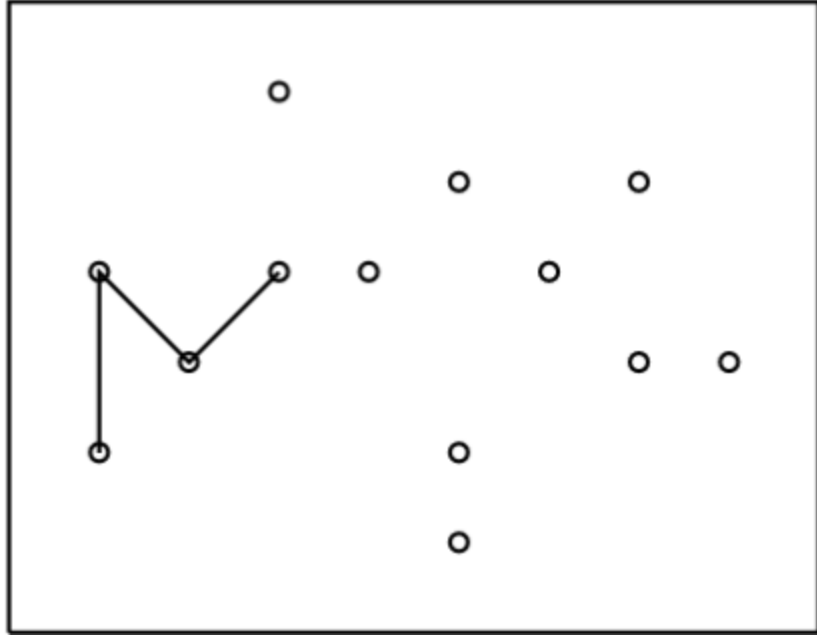


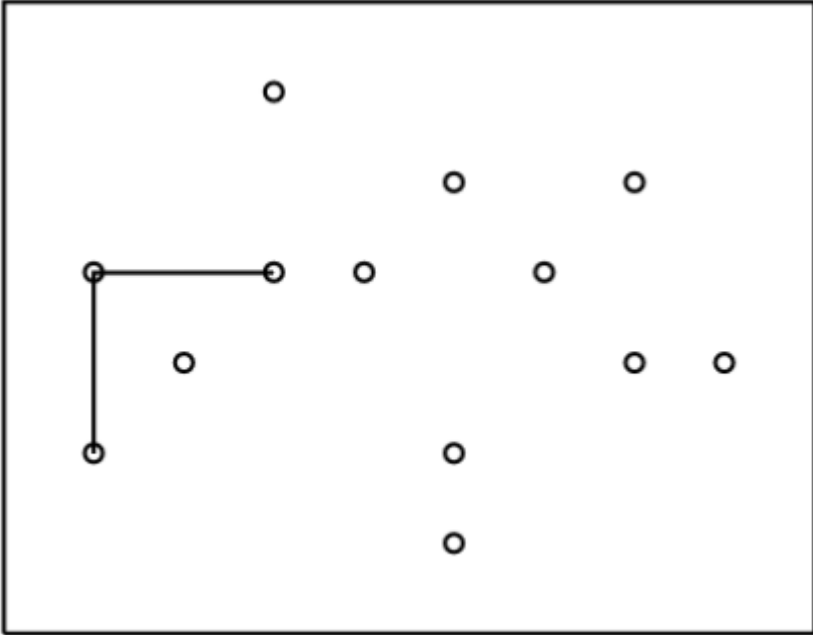
Andrew's algorithm

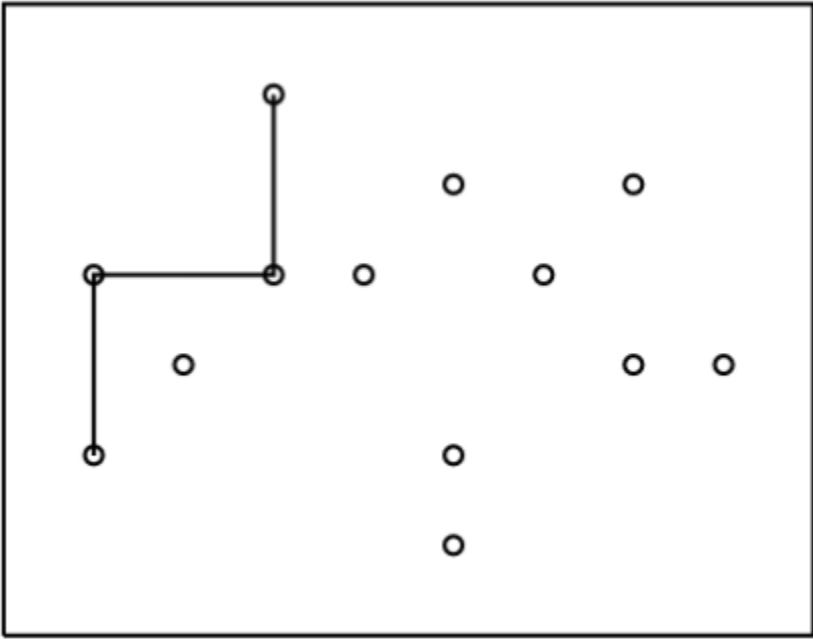
- The algorithm first determines the leftmost and rightmost points in the set, and then constructs the convex hull in two parts: first the upper hull and then the lower hull.
- Both parts are similar, so we can focus on constructing the upper hull.
- First, we sort the points primarily according to x coordinates and secondarily according to y coordinates

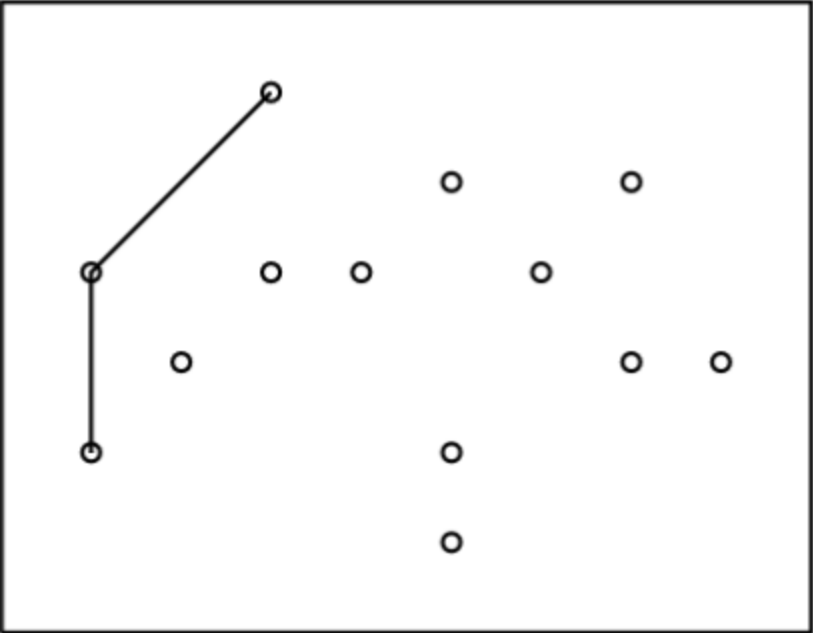


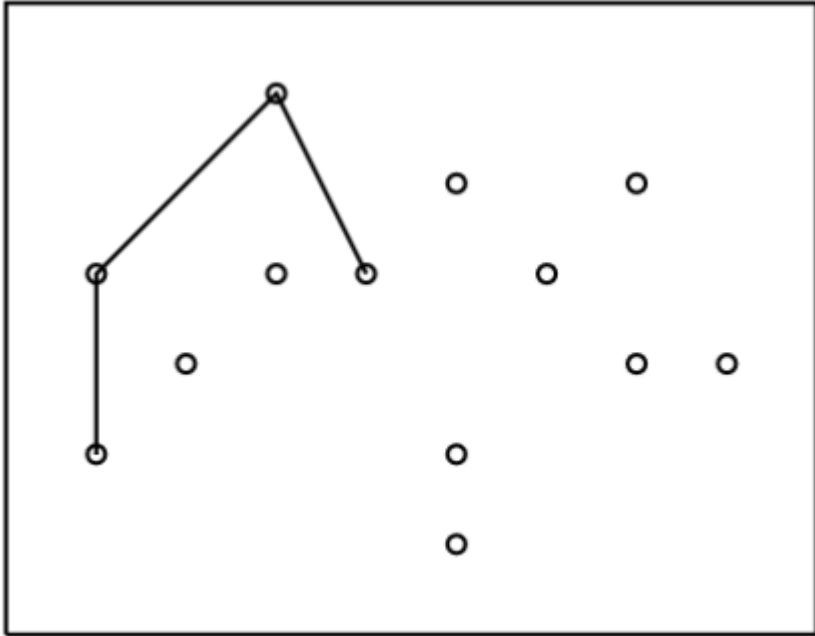


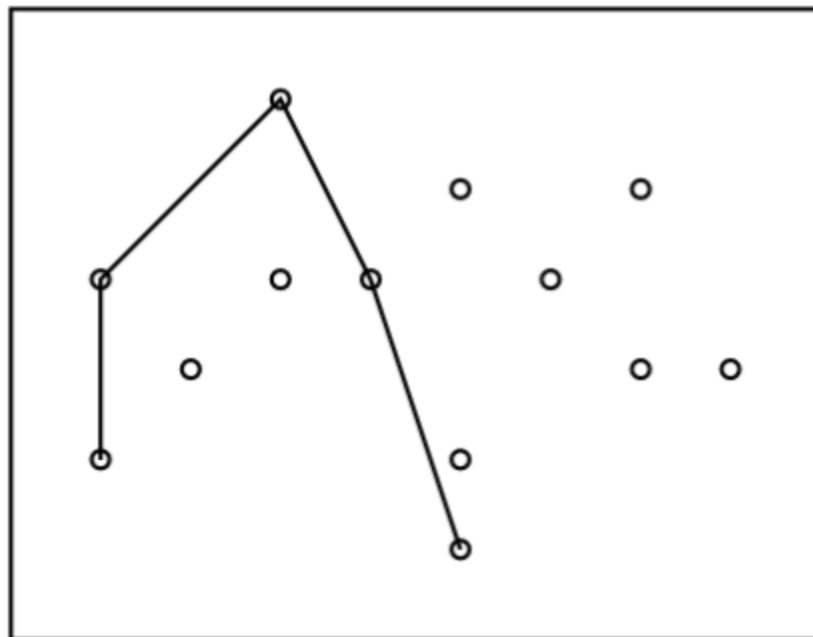


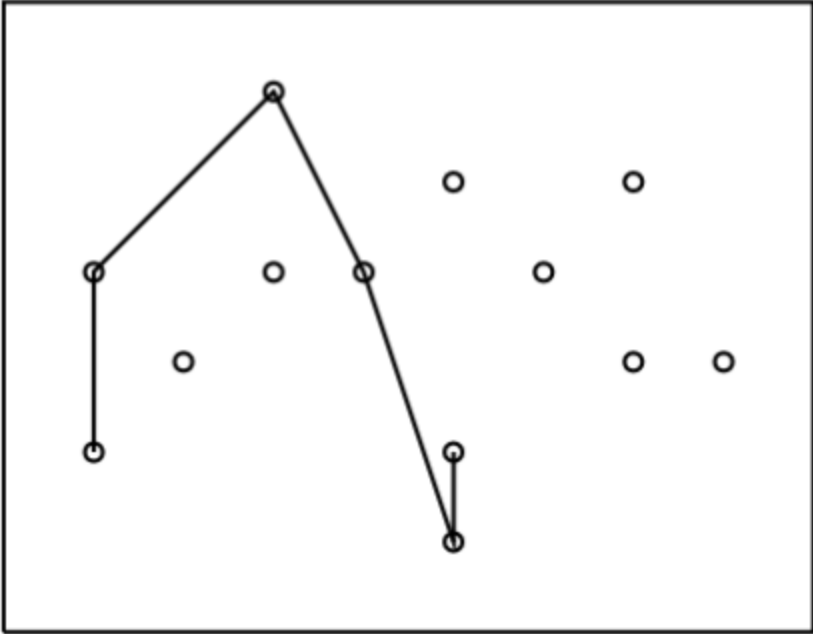


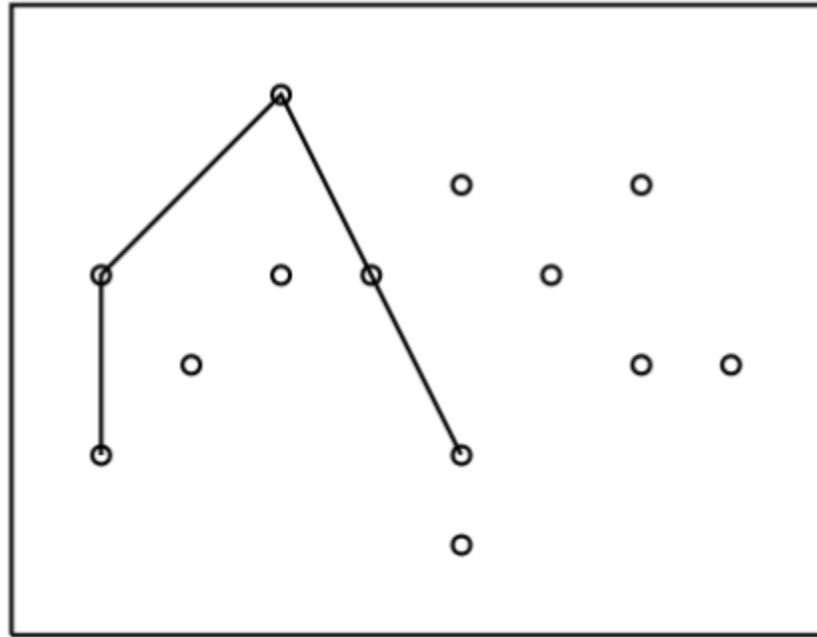


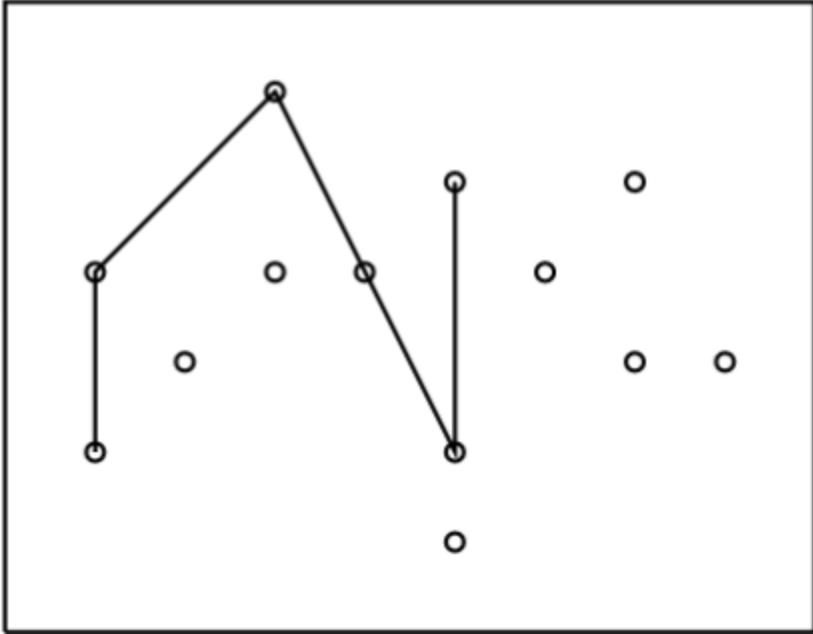


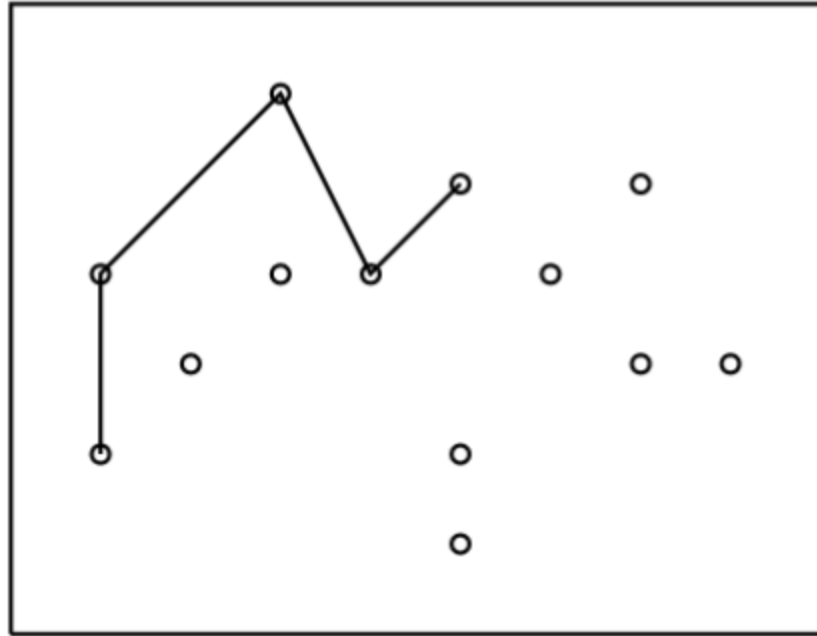


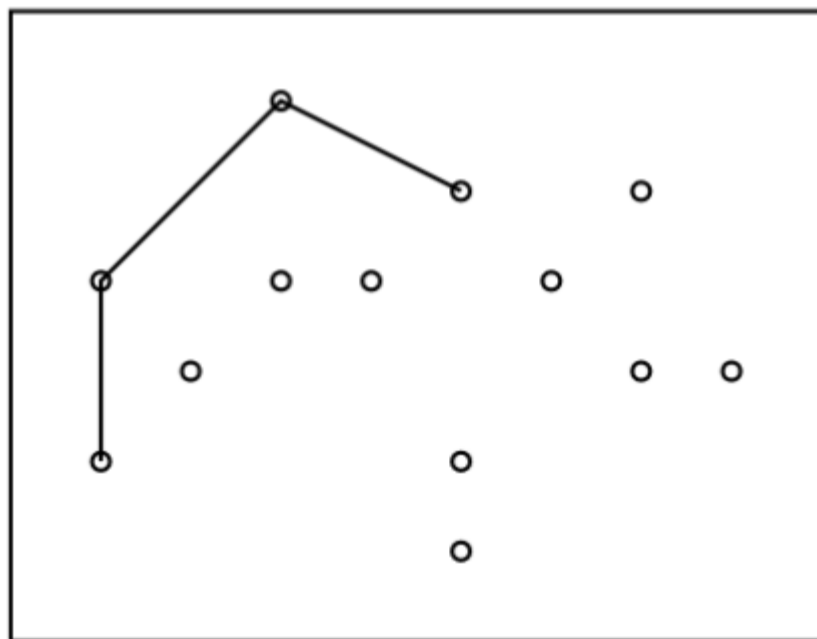


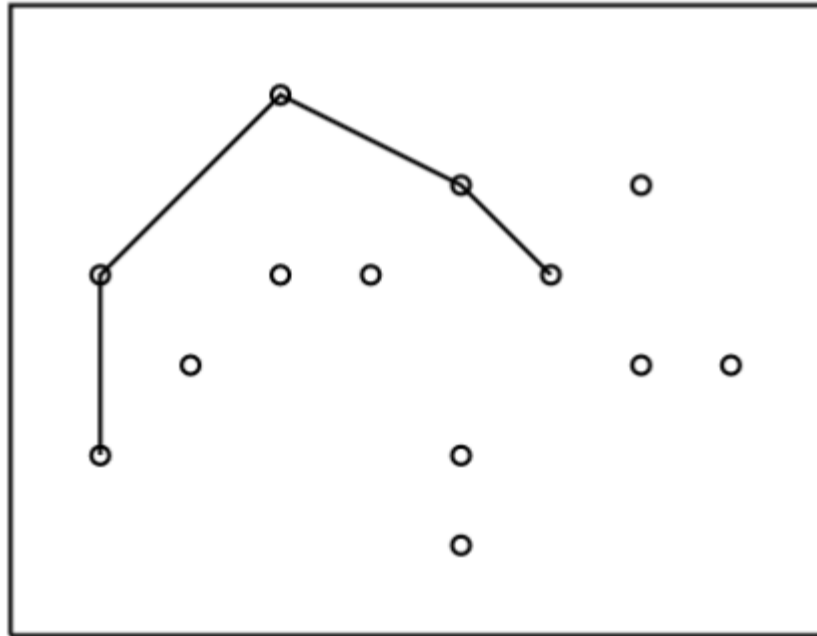


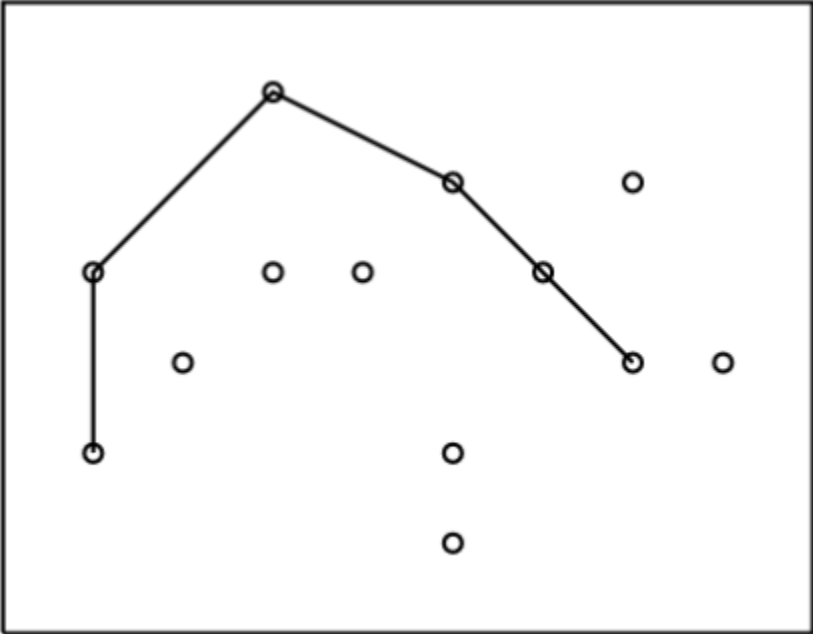


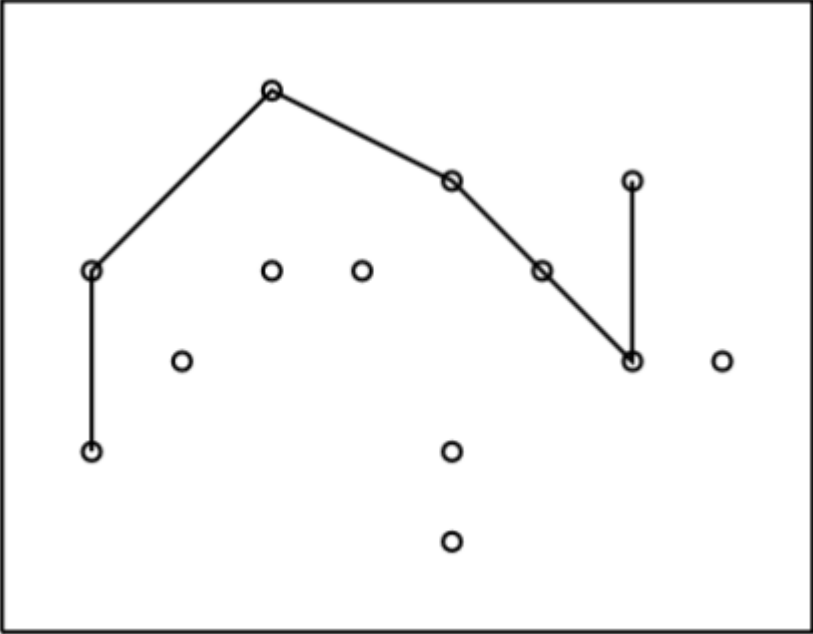


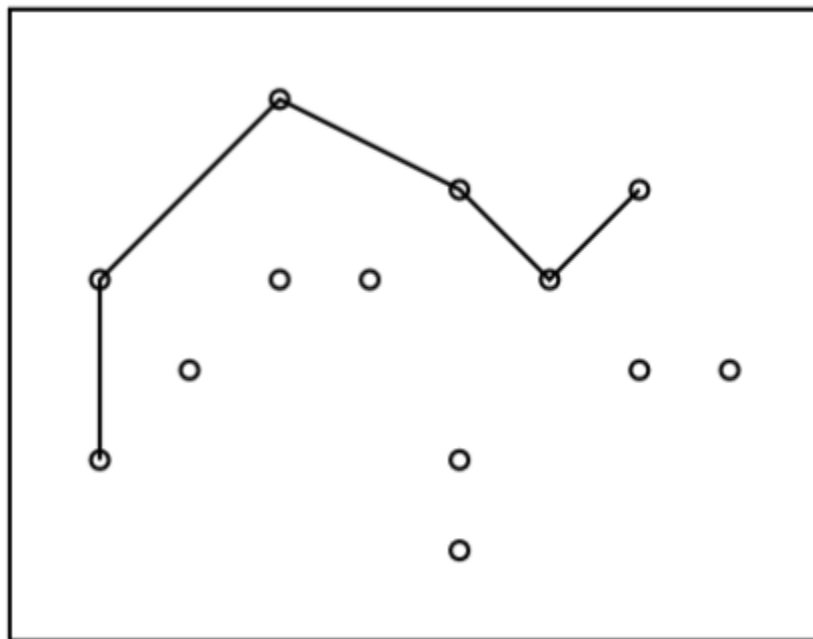


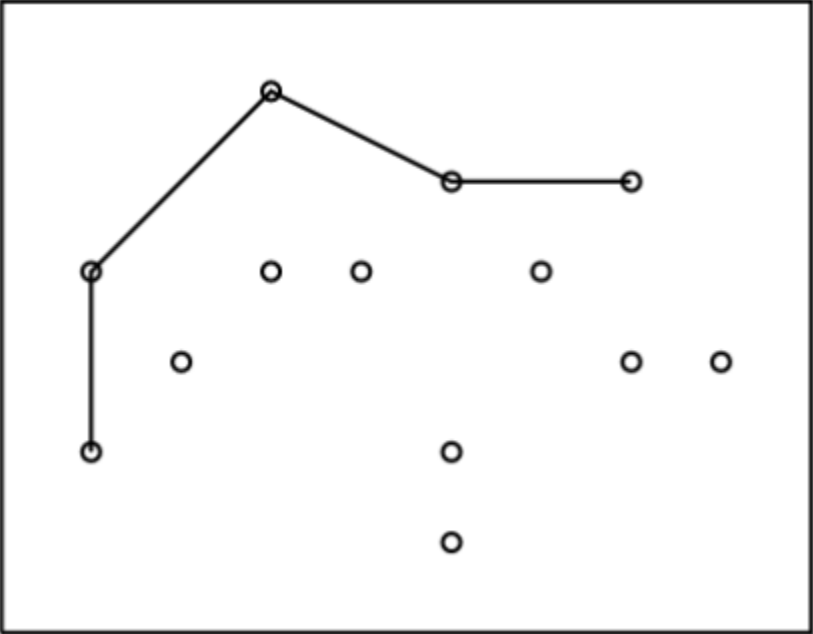


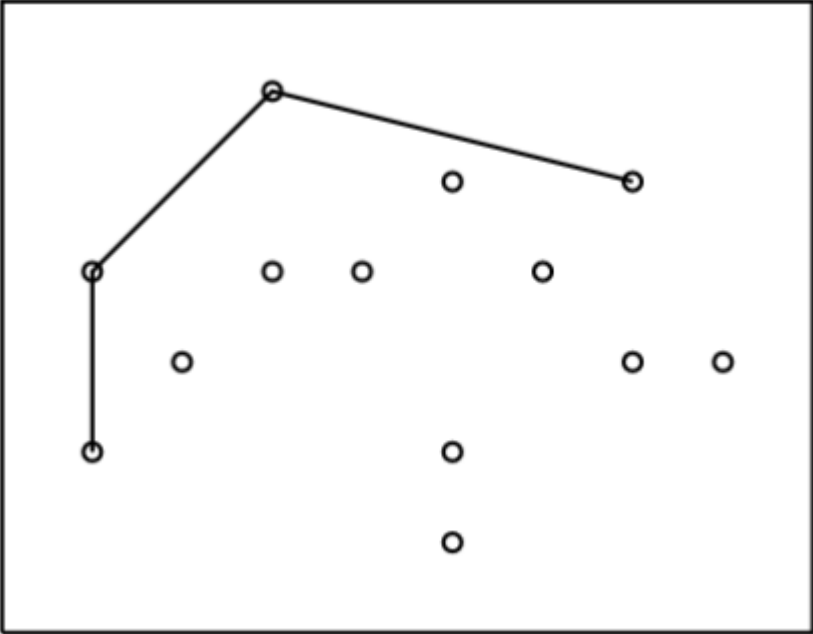


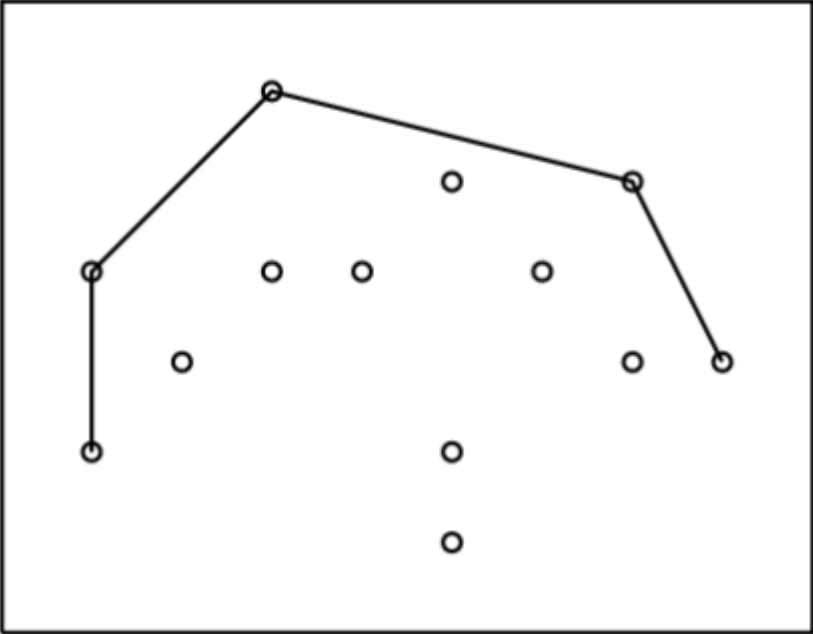


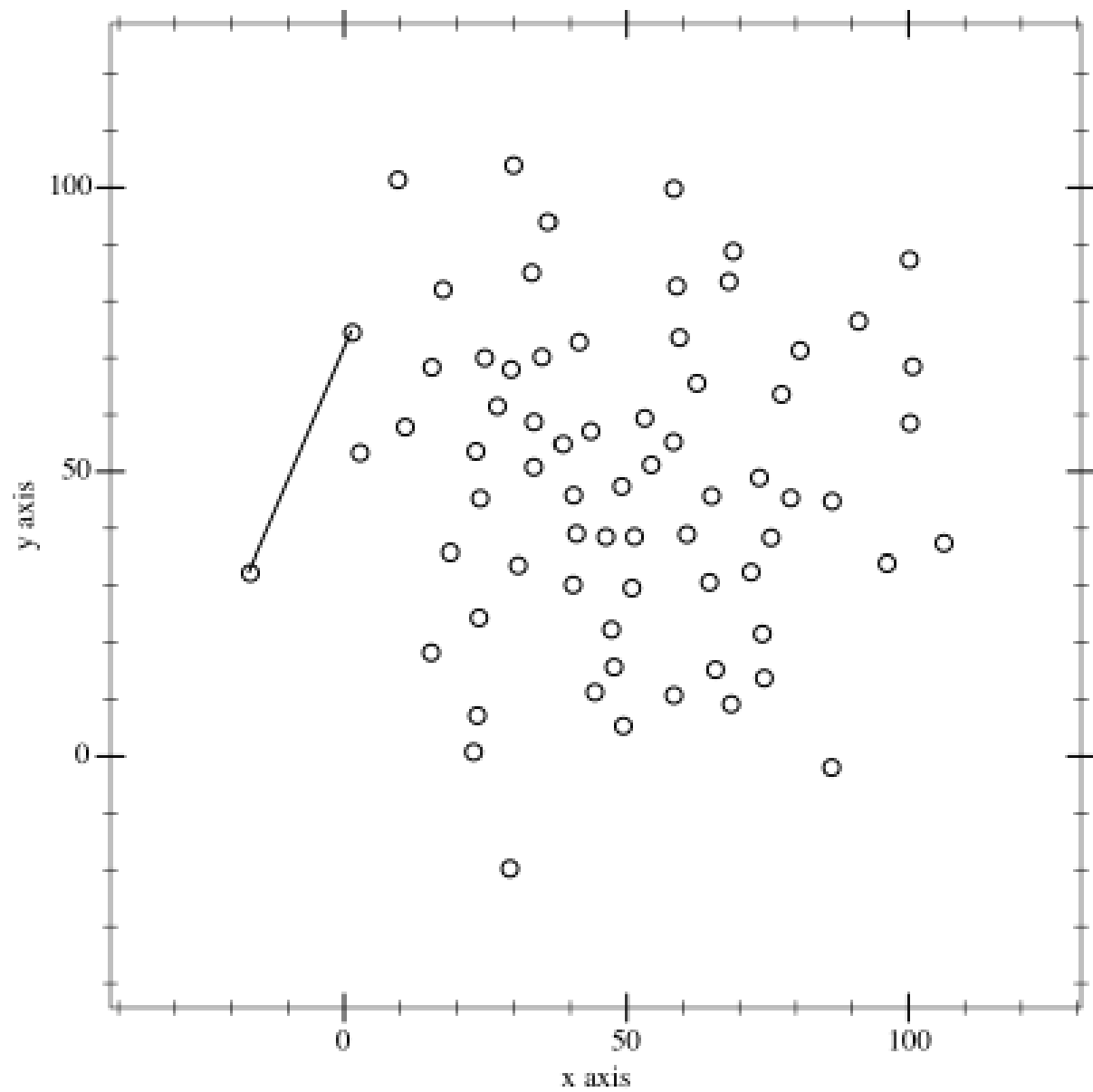












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